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CYCLIC PERMUTATIONS: CROSSING NUMBERS OF THE JOIN PRODUCTS OF GRAPHS

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Abstract. The crossing number cr(G) of a graph G is the minimal number of edge crossings over all drawings of G in the plane. In the paper, we extend known results concerning crossing numbers for join of graphs of order six. We give the crossing number of the join product $G + D_n$, where the graph G consists of one 4-cycle and two leaves, and D_n consists on n isolated vertices. The proof is done with the help of software that generates all cyclic permutations for a given number k, and creates a graph for a calculating the distances between all (k-1)! vertices of the graph. Finally, by adding some edges to the graph G, we are able to obtain the crossing numbers of the join product with the discrete graph D_n for other graphs.

Keywords: graph, drawing, crossing number, join product, cyclic permutation

Mathematics subject classification: Primary 05C10, 05C38

1 Introduction

Let G be a simple graph with the vertex set V and the edge set E. A drawing of the graph G is a representation of G in the plane such that its vertices are represented by distinct points and its edges by simple continuous arcs connecting the corresponding point pairs. In such a drawing, the intersection of the interiors of the arcs is called a *crossing*. A drawing is *good* if each two edges have at most one point in common, which is either a common end-vertex or a crossing. Moreover, no three edges cross in a point. It is easy to see that a drawing with minimum number of crossings (an optimal drawing) is always a good drawing. The *crossing number* cr(G) of a simple graph G is defined as the minimum possible number of edge crossings in a good drawing of G in the plane. Let G_1 and G_2 be simple graphs with vertex sets $V(G_1)$ and $V(G_2)$ and edge sets $E(G_1)$ and $E(G_2)$, respectively. The join product of two graphs G_1 and G_2 , denoted by $G_1 + G_2$, is obtained from the vertex-disjoint copies of G_1 and G_2 by adding all edges between $V(G_1)$ and $V(G_2)$. For $|V(G_1)| = m$ and $|V(G_2)| = n$, the edge set of $G_1 + G_2$ is the union of disjoint edge sets of the graphs G_1, G_2 , and the complete bipartite graph $K_{m,n}$. Let D(D(G)) be a good drawing of the graph G. We denote the number of crossings in D by $cr_D(G)$.

Let G_i and G_j be edge-disjoint subgraphs of G. We denote the number of crossings between edges of G_i and edges of G_j by $\operatorname{cr}_D(G_i, G_j)$, and the number of crossings among edges of G_i in D by $\operatorname{cr}_D(G_i)$. In the paper, some proofs are based on the Kleitman's result on crossing numbers of the complete bipartite graphs [12]. More precisely, he proved that

$$\operatorname{cr}(K_{m,n}) = \left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{m-1}{2} \right\rfloor \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor, \quad \text{if} \quad \min\{m,n\} \le 6.$$

2 The crossing number of $G + D_n$



Fig. 1. Drawing of the graph G with the vertex notation and the graph $G + D_2$.

In the paper, we extent these results [4], [6], [7], [9], [10], [11] for another four graphs. Let G be the graph consisting of one 4-cycle and of two leaves. We consider the join product of G with the discrete graph on n vertices denoted by D_n . The graph $G + D_n$ consists of one copy of the graph Gand of n vertices t_1, t_2, \ldots, t_n , where any vertex t_i , $i = 1, 2, \ldots, n$, is adjacent to every vertex of G. Let T^i , $1 \le i \le n$, denote the subgraph induced by the six edges incident with the vertex t_i . Thus, $T^1 \cup \cdots \cup T^n$ is isomorphic with the complete bipartite graph $K_{6,n}$ and

$$G + D_n = G \cup K_{6,n} = G \cup \left(\bigcup_{i=1}^n T^i\right).$$
(1)

2.1 Cyclic permutations and configurations

Let D be a good drawing of the graph $G + D_n$. The rotation $\operatorname{rot}_D(t_i)$ of a vertex t_i in the drawing D is the cyclic permutation that records the (cyclic) counter-clockwise order in which the edges leave t_i , see [3]. We emphasize that a rotation is a cyclic permutation. Hence, for $i, j \in \{1, 2, \ldots, n\}, i \neq j$, every subgraph $T^i \cup T^j$ of the graph $G + D_n$ is isomorphic with the graph $K_{6,2}$. In the paper, we will deal with the minimum necessary number of crossings between the edges of T^i and the edges of T^j in a subgraph $T^i \cup T^j$ induced by the drawing D of the graph $G + D_n$ depending on the rotations $\operatorname{rot}_D(t_i)$ and $\operatorname{rot}_D(t_j)$.

Let D be a good drawing of the graph $K_{6,n}$. Woodall [13] proved that in the subdrawing of $T^i \cup T^j \cong K_{6,2}$ induced by D, $\operatorname{cr}_D(T^i, T^j) \ge 6$ if $\operatorname{rot}_D(t_i) = \operatorname{rot}_D(t_j)$. Moreover, if $Q(\operatorname{rot}_D(t_i), \operatorname{rot}_D(t_j))$ denotes the minimum number of interchanges of adjacent elements of $\operatorname{rot}_D(t_i)$ required to produce the inverse cyclic permutation of $\operatorname{rot}_D(t_j)$, then $Q(\operatorname{rot}_D(t_i), \operatorname{rot}_D(t_j)) \le \operatorname{cr}_D(T^i, T^j)$. We will separate the subgraphs T^i , $i = 1, \ldots, n$, of the graph $G + D_n$ into three subsets depending on haw many the considered T^i crosses the edges of G in D. For $i = 1, 2, \ldots, n$, let $R_D = \{T^i : \operatorname{cr}_D(G, T^i) = 0\}$ and $S_D = \{T^i : \operatorname{cr}_D(G, T^i) = 1\}$. Every other subgraph T^i crosses G at least twice in D. Moreover, let F^i denote the subgraph $G \cup T^i$ for $T^i \in R_D$, where $i \in \{1, \ldots, n\}$. Thus, for a given drawing of

Name	Cyclic perm.	Name	Cyclic perm.	Name	Cyclic perm.	Name	Cyclic perm.
$P_1 \rightarrow$	(1 2 3 4 5 6)	$P_{31} \rightarrow$	(1 2 3 6 4 5)	$P_{61} \rightarrow$	(1 2 5 6 3 4)	$P_{91} \rightarrow$	(1 4 5 6 2 3)
$P_2 \rightarrow$	(1 3 2 4 5 6)	$P_{32} \rightarrow$	(1 3 2 6 4 5)	$P_{62} \rightarrow$	(152634)	$P_{92} \rightarrow$	(154623)
$P_3 \rightarrow$	(1 2 4 3 5 6)	$P_{33} \rightarrow$	(1 2 6 3 4 5)	$P_{63} \rightarrow$	(1 2 6 5 3 4)	$P_{93} \rightarrow$	(1 4 6 5 2 3)
$P_4 \rightarrow$	(1 4 2 3 5 6)	$P_{34} \rightarrow$	(1 6 2 3 4 5)	$P_{64} \rightarrow$	(162534)	$P_{94} \rightarrow$	(164523)
$P_5 \rightarrow$	(134256)	$P_{35} \rightarrow$	(1 3 6 2 4 5)	$P_{65} \rightarrow$	(156234)	$P_{95} \rightarrow$	(156423)
$P_6 \rightarrow$	(1 4 3 2 5 6)	$P_{36} \rightarrow$	(163245)	$P_{66} \rightarrow$	(165234)	$P_{96} \rightarrow$	(165423)
$P_7 \rightarrow$	(1 2 3 5 4 6)	$P_{37} \rightarrow$	(1 2 4 6 3 5)	$P_{67} \rightarrow$	(1 3 5 6 2 4)	$P_{97} \rightarrow$	(1 3 4 5 6 2)
$P_8 \rightarrow$	(1 3 2 5 4 6)	$P_{38} \rightarrow$	(1 4 2 6 3 5)	$P_{68} \rightarrow$	(153624)	$P_{98} \rightarrow$	(1 4 3 5 6 2)
$P_9 \rightarrow$	(1 2 5 3 4 6)	$P_{39} \rightarrow$	(1 2 6 4 3 5)	$P_{69} \rightarrow$	(1 3 6 5 2 4)	$P_{99} \rightarrow$	(1 3 5 4 6 2)
$P_{10} \rightarrow$	(152346)	$P_{40} \rightarrow$	(1 6 2 4 3 5)	$P_{70} \rightarrow$	(163524)	$P_{100} \rightarrow$	(1 5 3 4 6 2)
$P_{11} \rightarrow$	(1 3 5 2 4 6)	$P_{41} \rightarrow$	(1 4 6 2 3 5)	$P_{71} \rightarrow$	(156324)	$P_{101} \rightarrow$	(1 4 5 3 6 2)
$P_{12} \rightarrow$	(153246)	$P_{42} \rightarrow$	(164235)	$P_{72} \rightarrow$	(165324)	$P_{102} \rightarrow$	(154362)
$P_{13} \rightarrow$	(1 2 4 5 3 6)	$P_{43} \rightarrow$	(1 3 4 6 2 5)	$P_{73} \rightarrow$	(1 2 4 5 6 3)	$P_{103} \rightarrow$	(1 3 4 6 5 2)
$P_{14} \rightarrow$	(1 4 2 5 3 6)	$P_{44} \rightarrow$	(1 4 3 6 2 5)	$P_{74} \rightarrow$	(1 4 2 5 6 3)	$P_{104} \rightarrow$	(1 4 3 6 5 2)
$P_{15} \rightarrow$	(1 2 5 4 3 6)	$P_{45} \rightarrow$	(1 3 6 4 2 5)	$P_{75} \rightarrow$	(1 2 5 4 6 3)	$P_{105} \rightarrow$	(1 3 6 4 5 2)
$P_{16} \rightarrow$	(152436)	$P_{46} \rightarrow$	(163425)	$P_{76} \rightarrow$	(152463)	$P_{106} \rightarrow$	(163452)
$P_{17} \rightarrow$	(1 4 5 2 3 6)	$P_{47} \rightarrow$	(1 4 6 3 2 5)	$P_{77} \rightarrow$	(1 4 5 2 6 3)	$P_{107} \rightarrow$	(1 4 6 3 5 2)
$P_{18} \rightarrow$	(154236)	$P_{48} \rightarrow$	(164325)	$P_{78} \rightarrow$	(154263)	$P_{108} \rightarrow$	(164352)
$P_{19} \rightarrow$	(1 3 4 5 2 6)	$P_{49} \rightarrow$	(1 2 3 5 6 4)	$P_{79} \rightarrow$	(1 2 4 6 5 3)	$P_{109} \rightarrow$	(1 3 5 6 4 2)
$P_{20} \rightarrow$	(1 4 3 5 2 6)	$P_{50} \rightarrow$	(1 3 2 5 6 4)	$P_{80} \rightarrow$	(1 4 2 6 5 3)	$P_{110} \rightarrow$	(153642)
$P_{21} \rightarrow$	(1 3 5 4 2 6)	$P_{51} \rightarrow$	(1 2 5 3 6 4)	$P_{81} \rightarrow$	(1 2 6 4 5 3)	$P_{111} \rightarrow$	(1 3 6 5 4 2)
$P_{22} \rightarrow$	(153426)	$P_{52} \rightarrow$	(152364)	$P_{82} \rightarrow$	(1 6 2 4 5 3)	$P_{112} \rightarrow$	(163542)
$P_{23} \rightarrow$	(1 4 5 3 2 6)	$P_{53} \rightarrow$	(1 3 5 2 6 4)	$P_{83} \rightarrow$	(1 4 6 2 5 3)	$P_{113} \rightarrow$	(156342)
$P_{24} \rightarrow$	(154326)	$P_{54} \rightarrow$	(153264)	$P_{84} \rightarrow$	(164253)	$P_{114} \rightarrow$	(165342)
$P_{25} \rightarrow$	(1 2 3 4 6 5)	$P_{55} \rightarrow$	(1 2 3 6 5 4)	$P_{85} \rightarrow$	(1 2 5 6 4 3)	$P_{115} \rightarrow$	(1 4 5 6 3 2)
$P_{26} \rightarrow$	(1 3 2 4 6 5)	$P_{56} \rightarrow$	(1 3 2 6 5 4)	$P_{86} \rightarrow$	(152643)	$P_{116} \rightarrow$	(154632)
$P_{27} \rightarrow$	(1 2 4 3 6 5)	$P_{57} \rightarrow$	(1 2 6 3 5 4)	$P_{87} \rightarrow$	(1 2 6 5 4 3)	$P_{117} \rightarrow$	(1 4 6 5 3 2)
$P_{28} \rightarrow$	(1 4 2 3 6 5)	$P_{58} \rightarrow$	(1 6 2 3 5 4)	$P_{88} \rightarrow$	(1 6 2 5 4 3)	$P_{118} \rightarrow$	(164532)
$P_{29} \rightarrow$	(1 3 4 2 6 5)	$P_{59} \rightarrow$	(1 3 6 2 5 4)	$P_{89} \rightarrow$	(156243)	$P_{119} \rightarrow$	(156432)
$P_{30} \rightarrow$	(1 4 3 2 6 5)	$P_{60} \rightarrow$	(163254)	$P_{90} \rightarrow$	(165243)	$P_{120} \rightarrow$	(165432)

Tab. 1. Names of Cyclic Permutations of 6-elements set.

G, any F^i is exactly represented by $rot_D(t_i)$. All cyclic permutations of six elements can be generated by the algorithm [2], and they are collected in Tab. 1.

We will dealt with only drawings of the graph G with a possibility of an existence of a subgraph $T_i \in R_D$ because of arguments in the proof of the main Theorem 1. Assume a good drawing D of the graph $G + D_n$ in which the edges of G does not cross each other. In this case, without loss of generality, we can choose the vertex notations of the graph in such a way as shown in Fig. 1(a). It is easy to see that, in D, there are only four different possible configurations of F^i summarized in Tab. 2. In the rest of the paper, each cyclic permutation will be represented by the permutation with 1 in the first position. As for our considerations does not play role which of the regions is unbounded, assume the drawings shown in Figure 2. In a fixed drawing of the graph $G + D_n$, some configurations from the set $\mathcal{M} = \{A_1, A_2, A_3, A_4\}$ do not must appear. We denote by \mathcal{M}_D the set of all configurations that exist in the drawing D belonging to \mathcal{M} .



Fig. 2. Drawings of four possible configurations of graph F^i with the vertices of G denoted as in Fig. 1(a).

$A_1:(125643)$	$A_2:(132546)$
$A_3:(125463)$	$A_4:(132564)$

Tab. 2. Configurations of graph F^i with the vertices of G denoted as in Fig. 1(a).

—	A_1	A_2	A_3	A_4
A_1	6	4	5	5
A_2	4	6	5	5
A_3	5	5	6	5
A_4	5	5	5	6

Tab. 3. Lower-bounds of numbers of crossings for two configurations from \mathcal{M} .

Let X, Y be the configurations from \mathcal{M}_D . We shortly denote by $\operatorname{cr}_D(X, Y)$ the number of crossings in D between T^i and T^j for different $T^i, T^j \in R_D$ such that F^i, F^j have configurations X, Y, respectively. Finally, let $\operatorname{cr}(X, Y) = \min\{\operatorname{cr}_D(X, Y)\}$ over all good drawings of the graph $G + D_n$. In the next statements we are able to use the possibilities of the algorithm of the cyclic permutations of 6-elements set, see [2]. By $\overline{P_i}$ we will understand the inverse cyclic permutation to the permutation P_i , for $i = 1, \ldots, 120$. Woodall [13] defined the cyclic-ordered graph COG with the set of vertices $V = \{P_1, P_2, \ldots, P_{120}\}$, and with the set of edges E, where two vertices are joined by the edge if the vertices correspond to the permutations P_i and P_j , which are formed by the exchange of exactly two adjacent elements of the 6-tuple (i. e. an ordered set with 6 elements). Hence, if $d_{COG}(\text{"rot}_D(t_i)\text{", "rot}_D(t_j)\text{")}$ denotes the distance between two vertices correspond to the cyclic permutations $\operatorname{rot}_D(t_i)$ and $\operatorname{rot}_D(t_i)$ in the graph COG, then

$$d_{COG}(\operatorname{rot}_D(t_i), \operatorname{rot}_D(t_j)) = Q(\operatorname{rot}_D(t_i), \operatorname{rot}_D(t_j)) \leq \operatorname{cr}_D(T^i, T^j)$$

for any two different subgraphs T^i and T^j . The configurations A_1 and A_2 are represented by the cyclic permutations $P_{85} = (125643)$ and $P_8 = (132546)$, respectively. Using $\overline{P_8} = (164523) = P_{94}$ and $d_{COG}("P_{85}", "P_{94}") = 4$ we obtain $\operatorname{cr}(A_1, A_2) \ge 4$. The same reason gives $\operatorname{cr}(A_1, A_3) \ge 5$, $\operatorname{cr}(A_1, A_4) \ge 5$, $\operatorname{cr}(A_2, A_3) \ge 5$, $\operatorname{cr}(A_2, A_4) \ge 5$ and $\operatorname{cr}(A_3, A_4) \ge 4$. Moreover, by a discussion of possible subdrawings, we can verify that $\operatorname{cr}(A_3, A_4) \ge 5$. Clearly, also $\operatorname{cr}(A_k, A_k) \ge 6$ holds for any $k = 1, \ldots, 4$. Thus, all lower-bounds of number of crossing of configurations from \mathcal{M} are summarized in Tab. 3.

2.2 Main results

Lemma 1 Let D be a good drawing of $G + D_n$, n > 2, in which $\operatorname{cr}_D(T^i, T^j) \neq 0$ for any different subgraphs T^i and T^j . Let $2|R_D| + |S_D| > 2n - 2\left\lfloor \frac{n}{2} \right\rfloor$ and let $T^n, T^{n-1} \in R_D$ be different subgraphs with $\operatorname{cr}_D(T^n \cup T^{n-1}) \geq 4$. If both conditions

$$\operatorname{cr}_{D}(G \cup T^{n} \cup T^{n-1}, T^{i}) \ge 10 \qquad \qquad \text{for any } T^{i} \in R_{D} \setminus \{T^{n}, T^{n-1}\},$$

$$(2)$$

$$\operatorname{cr}_D(G \cup T^n \cup T^{n-1}, T^i) \ge 7 \qquad \qquad \text{for any } T^i \in S_D \tag{3}$$

hold, then there are at least $6\left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2\left\lfloor \frac{n}{2} \right\rfloor$ crossings in D.

*P*roof. We denote by $r = |R_D|$ and $s = |S_D|$. By the assumption of lemma, any $T^i \notin R_D \cup S_D$ satisfies the condition $\operatorname{cr}_D(G \cup T^n \cup T^{n-1}, T^i) \geq 4$, and the number of T^i that cross the graph G at least two times is equal to n - r - s. By fixing of the graph $G \cup T^n \cup T^{n-1}$ we have

$$\operatorname{cr}_{D}(G+D_{n}) = \operatorname{cr}_{D}(K_{6,n-2}) + \operatorname{cr}_{D}(K_{6,n-2}, G \cup T^{n} \cup T^{n-1}) + \operatorname{cr}_{D}(G \cup T^{n} \cup T^{n-1}) \geq \\ \geq 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 10(r-2) + 7s + 4(n-r-s) + 4 = 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 6r + 3s + 4n - 16 \geq \\ \geq 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 3 \left(2n-2 \left\lfloor \frac{n}{2} \right\rfloor + 1 \right) + 4n - 16 \geq 6 \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2 \left\lfloor \frac{n}{2} \right\rfloor.$$
This completes the proof

This completes the proof.

Lemma 2 Let D be a good drawing of $G+D_n$ with the vertex notations of the graph G as in Fig. 1(a), n > 2. If $T^n \in R_D$ such that F^n has configuration $A_i \in \mathcal{M}_D$, for i = 1, 3, 4, then

$$\operatorname{cr}_D(T^n, T^k) \ge 3$$
 for any $T^k \in S_D$. (4)

Proof. Let, in D, the graph F^n has configuration A_1 . If $T^k \in S_D$ with $\operatorname{cr}_D(T^n, T^k) = 2$, then the vertex t_k must be placed in a region with at least three vertices of G on its boundary, see Fig. 2. Since $T^k \in S_D$, the vertex t_k cannot be placed in the region bounded by 4-cycle of the graph G. Moreover, if t_k is placed in another regions, then $cr_D(F^n, T^k) > 3$. The same idea can be used for configurations A_3 and A_4 . This completes the proof.

Remark that the property (4) is not true for configuration A_2 , see the proof of the following statement.

Collorary 1 Let D be a good drawing of $G + D_n$ with the vertex notations of the graph G as in Fig. 1(a), n > 2, in which $\operatorname{cr}_D(T^i, T^j) \neq 0$ for any different subgraphs T^i and T^j . If $T^n, T^{n-1} \in R_D$ such that F^n , F^{n-1} have configurations A_1 , A_2 , respectively, then

$$\operatorname{cr}_D(G \cup T^n \cup T^{n-1}, T^k) \ge 7 \qquad \qquad \text{for any } T^k \in S_D. \tag{5}$$

Proof. Let, in D, the graphs F^n , F^{n-1} have configurations A_1, A_2 , respectively. The configurations A_1 and A_2 are represented by the cyclic permutations $P_{85} = (125643)$ and $P_8 = (132546)$, respectively.

• If there is a subgraph $T^k \in S_D$ with $\operatorname{cr}_D(T^{n-1}, T^k) = 2$, then the vertex t_k must be placed in the region with four vertices of G and one vertex t_{n-1} on its boundary, see Fig. 2. Thus, the graph $F^k = G \cup T^k$ can be represented only by two possible cyclic permutations $P_{81} =$ (126453) and $P_{95} = (156423)$. By the above mentioned algorithm we have

$$d_{COG}("P_{26}", "P_{85}") = d_{COG}("P_{99}", "P_{85}") = 4,$$

where $\overline{P_{81}} = (135462) = P_{99}$ and $\overline{P_{95}} = (132465) = P_{26}$. By the properties of the cyclic permutations we have $cr_D(T^n, T^k) \ge 4$. Thus, $cr_D(G \cup T^n \cup T^{n-1}, T^k) \ge 1 + 4 + 2 = 7$.

• If $\operatorname{cr}_D(T^{n-1}, T^k) \ge 3$ for any subgraph $T^k \in S_D$, then $\operatorname{cr}_D(G \cup T^n \cup T^{n-1}, T^k) \ge 1 + 3 + 3 = 7$.



Fig. 3. Two good drawings of $G + D_n$.

Theorem 1 $\operatorname{cr}(G+D_n) = 6\left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2\left\lfloor \frac{n}{2} \right\rfloor \text{ for } n \ge 1.$

Proof. In Fig. 3 there are the drawings of $G + D_n$ with $6 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + 2 \lfloor \frac{n}{2} \rfloor$ crossings. Thus, $\operatorname{cr}(G + D_n) \leq 6 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + 2 \lfloor \frac{n}{2} \rfloor$. We prove the reverse inequality by induction on n. The graph $G + D_1$ is planar, hence $\operatorname{cr}(G + D_1) = 0$. It is clear from Fig. 1(b) that $\operatorname{cr}(G + D_2) \leq 2$. The graph $G + D_2$ contains a subdivision of $K_{3,4}$, and therefore $\operatorname{cr}(G + D_2) \geq 2$. So, $\operatorname{cr}(G + D_2) = 2$ and the result is true for n = 1 and n = 2.

Suppose now that, for $n \ge 3$, there is a drawing D with

$$\operatorname{cr}_{D}(G+D_{n}) < 6\left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2\left\lfloor \frac{n}{2} \right\rfloor, \tag{6}$$

and let

$$\operatorname{cr}(G+D_m) \ge 6\left\lfloor \frac{m}{2} \right\rfloor \left\lfloor \frac{m-1}{2} \right\rfloor + 2\left\lfloor \frac{m}{2} \right\rfloor \qquad \text{for any } m < n. \tag{7}$$

The drawing D has the following property:

$$\operatorname{cr}_D(T^i, T^j) \neq 0 \qquad \text{for all } i, j = 1, 2, \dots, n, \ i \neq j.$$
(8)

To prove it assume that there are two different subgraphs T^i and T^j such that $\operatorname{cr}_D(T^i, T^j) = 0$. Without loss of generality let $\operatorname{cr}_D(T^{n-1}, T^n) = 0$. One can easy to verify that $\operatorname{cr}_D(G, T^{n-1} \cup T^n) \ge 2$. As $cr(K_{6,3}) = 6$, we have $cr_D(T^k, T^{n-1} \cup T^n) \ge 6$ for k = 1, 2, ..., n-2. So, for the number of crossings in D holds

$$\operatorname{cr}_{D}(G+D_{n}) = \operatorname{cr}_{D}(G+D_{n-2}) + \operatorname{cr}_{D}(T^{n-1}\cup T^{n}) + \operatorname{cr}_{D}(K_{6,n-2}, T^{n-1}\cup T^{n}) + \operatorname{cr}_{D}(G, T^{n-1}\cup T^{n}) \ge 6\left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 2\left\lfloor \frac{n-2}{2} \right\rfloor + 6(n-2) + 2 = 6\left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2\left\lfloor \frac{n}{2} \right\rfloor.$$

This contradicts (6), and therefore $\operatorname{cr}_D(T^i, T^j) \neq 0$ for all $i, j = 1, 2, \ldots, n, i \neq j$. Our assumption on D together with $\operatorname{cr}(K_{6,n}) = 6 \left| \frac{n}{2} \right| \left| \frac{n-1}{2} \right|$ implies that

$$\operatorname{cr}_D(G) + \operatorname{cr}_D(G, K_{6,n}) < 2 \left\lfloor \frac{n}{2} \right\rfloor.$$

Hence, if we denote $r = |R_D|$ and $s = |S_D|$, then

$$0r + 1s + 2(n - r - s) < 2\left\lfloor \frac{n}{2} \right\rfloor.$$

Thus, $r \ge 1$ and $2r + s > 2n - 2\lfloor \frac{n}{2} \rfloor$. We will fix one or two subgraphs with a contradiction with the assumption that there are less than $6\lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + 2\lfloor \frac{n}{2} \rfloor$ crossings in the following cases:

Case 1:
$$\operatorname{cr}_D(G) = 0$$
.

We will deal with the sets of configurations $\{A_1, A_2\}$ in the drawing D.

- 1) $\{A_1, A_2\} \not\subseteq \mathcal{M}_D$.
 - a) Let A₂ ∉ M_D and A_i ∈ M_D for some i ∈ {1,3,4}, or let A₂ ∈ M_D and A_i ∈ M_D for some i ∈ {3,4}. Without lost of generality, we can assume that Tⁿ ∈ R_D with Fⁿ having configuration A_i. Thus, by fixing of the graph Fⁿ using Lemma 2 we have

$$\operatorname{cr}_{D}(G+D_{n}) = \operatorname{cr}_{D}(K_{6,n-1}) + \operatorname{cr}_{D}(K_{6,n-1}, G \cup T^{n}) + \operatorname{cr}_{D}(G \cup T^{n}) \ge 6\left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor + 5(r-1) + 4s + 3(n-r-s) = 6\left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor + 2r + s + 3n - 5 \ge 6\left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor + 2n - 2\left\lfloor \frac{n}{2} \right\rfloor + 1 + 3n - 5 \ge 6\left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2\left\lfloor \frac{n}{2} \right\rfloor.$$

b) Let $\mathcal{M}_D = \{A_2\}$ and, without lost of generality, let $T^n \in R_D$.

If there is no $T^k \in S_D$ with $\operatorname{cr}_D(T^n, T^k) = 2$, then we fix the graph F^n having configuration A_2 and we obtain the same inequalities as in the previous case. So, assume that there is a subgraph $T^k \in S_D$ with $\operatorname{cr}_D(T^n, T^k) = 2$. We can easily verify that $\operatorname{cr}_D(G \cup T^n \cup T^k, T^i) \ge 6 + 2 = 8$ for any $T^i \in R_D$, because both F^n and F^i have configuration A_2 . Similarly by a discussion for two possible drawings of the graph T^k , see the proof of Corollary 1, we can verify that $\operatorname{cr}_D(G \cup T^n \cup T^k, T^i) \ge 7$ for any $T^i \in S_D$ and $\operatorname{cr}_D(G \cup T^n \cup T^k, T^i) \ge 6$ for any $T^i \notin R_D \cup S_D$. Thus, by fixing of the graph $G \cup T^n \cup T^k$ we have

$$\operatorname{cr}_{D}(G+D_{n}) = \operatorname{cr}_{D}(K_{6,n-2}) + \operatorname{cr}_{D}(K_{6,n-2}, G \cup T^{n} \cup T^{k}) + \operatorname{cr}_{D}(G \cup T^{n} \cup T^{k}) \geq \\ \geq 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 8(r-1) + 7s + 6(n-r-s) + 3 = 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 2r + s + \\ + 6n - 12 \geq 6 \left\lfloor \frac{n-2}{2} \right\rfloor \left\lfloor \frac{n-3}{2} \right\rfloor + 2n - 2 \left\lfloor \frac{n}{2} \right\rfloor + 1 + 6n - 12 \geq 6 \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2 \left\lfloor \frac{n}{2} \right\rfloor.$$

2) $\{A_1, A_2\} \subseteq \mathcal{M}_D$.

Without lost of generality let us fix any two T^n , $T^{n-1} \in R_D$ such that F^n , F^{n-1} have configurations A_1, A_2 , respectively. Then condition (2) is true by Tab. 3 and condition (3) holds by Corollary 1. Thus, all assumption of Lemma 1 are fulfilled.

Case 2: $cr_D(G) = 1$.



Fig. 4. Four possible drawings of the graph G with one crossing among its edges.

Since $r \ge 1$, without lost of generality we assume $T^n \in R_D$. In all four possible drawing of the graph G it is possible to verify that $\operatorname{cr}_D(G \cup T^n, T^i) \ge 4$ for any subgraph T^i , $i = 1, \ldots, n-1$. Thus, by fixing of the graph F^n we obtain

$$\operatorname{cr}_{D}(G+D_{n}) = \operatorname{cr}_{D}(K_{6,n-1}) + \operatorname{cr}_{D}(K_{6,n-1}, G \cup T^{n}) + \operatorname{cr}_{D}(G \cup T^{n}) \ge$$
$$\geq 6 \left\lfloor \frac{n-1}{2} \right\rfloor \left\lfloor \frac{n-2}{2} \right\rfloor + 4(n-1) + 1 \ge 6 \left\lfloor \frac{n}{2} \right\rfloor \left\lfloor \frac{n-1}{2} \right\rfloor + 2 \left\lfloor \frac{n}{2} \right\rfloor.$$

Case 3: $\operatorname{cr}_D(G) \ge 2$.

We are able to use the same idea as in the previous case for all possible drawing of the graph G with a possibility of an existence of a subgraph $T^i \in R_D$ in the considering drawing D.

This completes the proof of the main theorem.

2.3 Corollaries



Fig. 5. Four graphs G_1 , G_2 , G_3 , and G_4 by adding new edges to the graph G.

In Fig. 2 we are able to add some edges to the graph G without another crossings. So the drawing of the graphs $G_1 + D_n$, $G_2 + D_n$, $G_3 + D_n$, and $G_4 + D_n$ with $6\lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + 2\lfloor \frac{n}{2} \rfloor$ crossings is obtained. Thus, the next results are obvious.

Collorary 2 cr
$$(G_i + D_n) = 6 \lfloor \frac{n}{2} \rfloor \lfloor \frac{n-1}{2} \rfloor + 2 \lfloor \frac{n}{2} \rfloor$$
 for $n \ge 1$, where $i = 1, \dots, 4$.

Remark that the crossing numbers of the graphs $G_3 + D_n$ and $G_4 + D_n$ were obtained in [8], [5] without using the vertex rotation.

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